

# Concavity and random LPs

(1)

- In the process of bounding the revenue in the single-resource allocation problem, we obtained the following LP

$$V_n^{UB}(c | \{D_1, D_2, \dots, D_n\}) \equiv \max \sum_{i=1}^n p_i x_i$$

↑  
upper bound given random demands  $D_1, D_2, \dots, D_n$

$$\text{s.t. } \sum_{i=1}^n x_i \leq c$$
$$x_i \leq D_i \quad \forall i$$
$$x_i \geq 0 \quad \forall i$$

This is sometimes called the randomized-LP bound

- We also defined the fluid-LP bound

$$V_n^{FL}(c) \equiv \max \sum_{i=1}^n p_i x_i$$
$$\text{s.t. } x_i \leq \mu_i \quad (\mu_i = E[D_i])$$
$$\sum_{i=1}^n x_i \leq c$$
$$x_i \geq 0 \quad \forall i$$

②

Finally we showed

$$V_n^{Fl}(c) \geq E[V_n^{UB}(c | \{D_i\})]$$

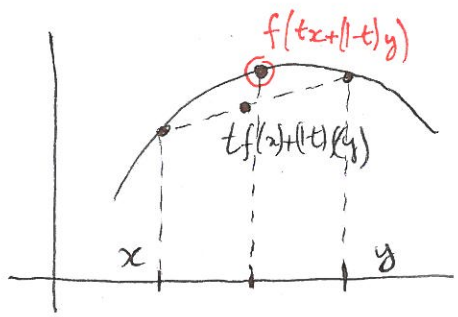
We will now see how to derive such results directly using convexity and Jensen's Inequality.

Main tools from convexity

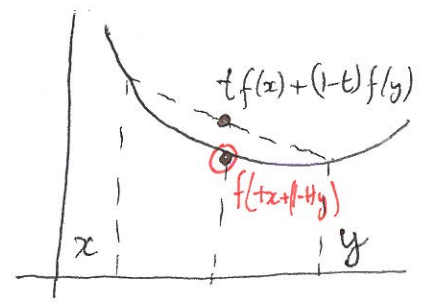
• Defn - A function  $f(x)$  is convex if  $\forall x, y$  and  $\forall t \in [0, 1]$ , the function obeys

$$(*) \quad f(tx + (1-t)y) \leq tf(x) + (1-t)f(y)$$

- Notes - i) Function  $f(x)$  is concave if  $-f(x)$  is convex (alternately, replace  $\leq$  with  $\geq$  in  $(*)$  to get concavity)
- ii) The definition works even if  $x$  is a vector
- iii) Graphical way to remember -



- Suppose  $t = 1/2$
- $tx + (1-t)y$  is the pt midway bet<sup>n</sup>  $x, y$



• Operations which preserve convexity (see HW2) <sup>(3)</sup>

i) Scaling -  $f(x)$  convex,  $a \geq 0$ , then

$f(ax+b)$  is convex

ii) ~~Linear~~ <sup>Linear</sup> Combination -  $f_1(x), f_2(x)$  convex,  $a_1, a_2 \geq 0$

$\Rightarrow a_1 f_1(x) + a_2 f_2(x)$  is convex

iii) Maximization -  $f_1(x), f_2(x)$  convex, then

$\max(f_1(x), f_2(x))$  is convex

(there are many others, but these suffice for us)

• For concavity, it is preserved under scaling, linear combination and minimization.

• A linear function  $f(x) = ax + b$  is both convex and concave.

• Jensen's Inequality - If  $f(x)$  is ~~convex~~ <sup>convex</sup>, then

$$E[f(x)] \geq f(E[x])$$

(  $\leq$  concave )

Now we show how to use these for LPs <sup>(4)</sup>

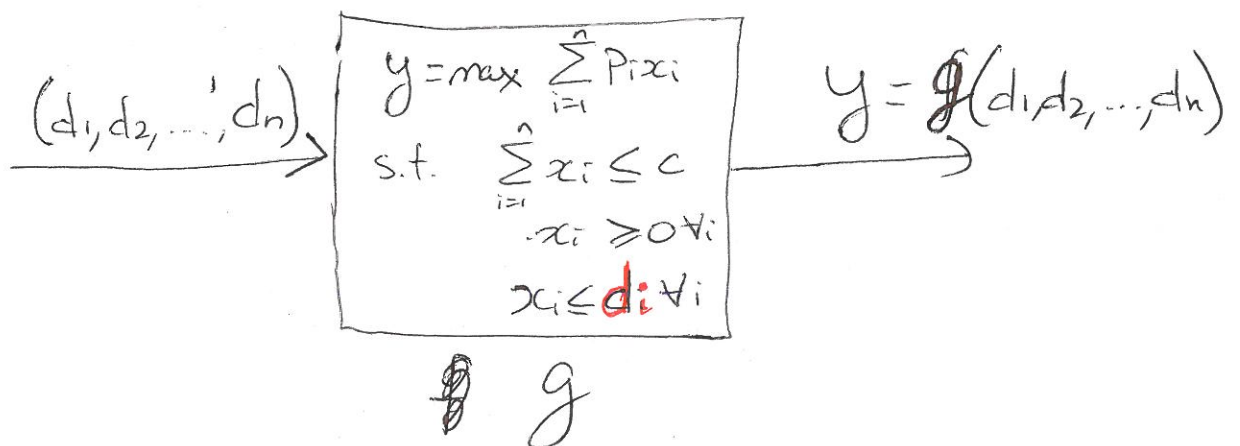
## • Primal Argument

- We want to show that the expected value of the randomized LP is bounded by the fluid LP.

$$\mathbb{E} \left[ \begin{array}{l} \max \sum_{i=1}^n p_i x_i \\ \text{s.t.} \sum_{i=1}^n x_i \leq c \\ x_i \leq D_i \quad \forall i \\ x_i \geq 0 \quad \forall i \end{array} \right] \leq \left[ \begin{array}{l} \max \sum_{i=1}^n p_i x_i \\ \text{s.t.} \sum_{i=1}^n x_i \leq c \\ x_i \leq \mathbb{E}[D_i] \quad \forall i \\ x_i \geq 0 \end{array} \right]$$

**\*\***

- To see that this can be shown via Jensen's, we need to first understand that the optimization problem above is a function  $f$  which maps 'capacities'  $(D_1, D_2, \dots, D_n)$  to an output. Visualize this as follows



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Thus, if we show  $g(d_1, d_2, \dots, d_n)$  is ~~convex~~ concave in  $(d_1, d_2, \dots, d_n)$ , we can use Jensen's Inequality to get **(\*\*)**

- Suppose we have 2 inputs  $\underline{d} = (d_1, d_2, \dots, d_n)$   
 $\underline{\delta} = (\delta_1, \delta_2, \dots, \delta_n)$

$$\text{let } \underline{x}_d^* \equiv \underset{\substack{\text{s.t. } \sum_{i=1}^n x_i \leq c \\ 0 \leq x_i \leq d_i \forall i}}{\text{argmax}} \sum_{i=1}^n p_i x_i \quad \left| \quad \underline{x}_s^* \equiv \underset{\substack{\text{s.t. } \sum_{i=1}^n x_i \leq c \\ 0 \leq x_i \leq \delta_i}}{\text{argmax}} \sum_{i=1}^n p_i x_i$$

(and  $g(\underline{d})$ ,  $g(\underline{\delta})$  the corresponding objective values)

- For some  $t \in [0, 1]$ , consider  $\underline{\Delta} = t\underline{d} + (1-t)\underline{\delta}$   
(i.e., for each  $i$ ,  $\Delta_i = td_i + (1-t)\delta_i$ )

To show ~~convexity~~ concavity, we need to show

$$g(\underline{\Delta}) \geq t g(\underline{d}) + (1-t) g(\underline{\delta})$$

To see this, we first observe that ~~the~~ the point  $x_{\Delta} = t x_d^* + (1-t) x_s^*$  is feasible for demand vector  $\Delta$ , i.e

$$i) \sum_{i=1}^n x_{\Delta,i} = t \sum_{i=1}^n x_{d,i}^* + (1-t) \sum_{i=1}^n x_{s,i}^* \leq t c + (1-t) c = c$$

$$ii) x_{\Delta,i} = t x_{d,i}^* + (1-t) x_{s,i}^* \leq t d_i + (1-t) s_i = \Delta_i$$

$$\Rightarrow \sum_{i=1}^n P_i x_{\Delta,i} \leq \left[ \begin{array}{l} \max_x \sum P_i x_i \\ \text{s.t. } \sum x_i \leq c \\ 0 \leq x_i \leq \Delta_i \forall i \end{array} \right] = g(\underline{\Delta})$$

$$\text{However } \sum_{i=1}^n P_i x_{\Delta,i} = t \sum_{i=1}^n P_i x_{d,i}^* + (1-t) \sum_{i=1}^n P_i x_{s,i}^* = t g(\underline{d}) + (1-t) g(\underline{s})$$

Thus  $g(\underline{\Delta}) \geq t g(\underline{d}) + (1-t) g(\underline{s}) \Rightarrow g$  is concave.

# Dual argument

We will now see an alternate (easier?) way to see the same result via duality

- We need one fact for this: Earlier, we said that if  $f_1, f_2$  are convex, then  $\max(f_1(x), f_2(x))$  is convex. Similarly, if  $g_1, g_2$  are concave fns, then  $\min(g_1(x), g_2(x))$  is concave. (try to check this!)
- Now, by LP duality, we have

$$g(d_1, \dots, d_n) = \left[ \begin{array}{l} \max \sum_{i=1}^n p_i x_i \\ \text{s.t. } \sum_{i=1}^n x_i \leq c \\ 0 \leq x_i \leq d_i \quad \forall i \end{array} \right] \begin{array}{l} \text{(dual)} \\ \text{vars} \\ : z \\ : \beta_i \end{array} \quad \underline{\text{PRIMAL}}$$

$$= \left[ \begin{array}{l} \min \quad cz + \sum_{i=1}^n d_i \beta_i \\ \text{s.t. } \beta_i + z \geq p_i \quad \forall i \\ \beta_i \geq 0, z \geq 0 \end{array} \right]$$

• Now let  $\mathcal{X}$  be the set of extreme points of the dual constraint set, i.e.

$$\mathcal{X} = \left\{ (\underline{\beta}, z) \mid (\underline{\beta}, z) \text{ is an extreme pt of } \begin{cases} \beta_i + z \geq p_i \forall i \\ \beta_i \geq 0, z \geq 0 \end{cases} \right\}$$

Then 
$$g(d_1, \dots, d_n) = \min_{(\underline{\beta}, z) \in \mathcal{X}} \left\{ \underbrace{Cz + \sum_{i=1}^n d_i \beta_i}_{\text{linear (hence concave) in } (d_1, d_2, \dots, d_n)} \right\}$$

$\Rightarrow g(d_1, \dots, d_n)$  is concave in  $d_1, \dots, d_n$

Thus - 
$$\mathbb{E} [g(D_1, D_2, \dots, D_n)] \leq g(\mathbb{E}[D_1], \dots, \mathbb{E}[D_n])$$

$$\mathbb{E} \left[ \begin{array}{l} \max \sum_{i=1}^n x_i p_i \\ \text{s.t. } \sum x_i \leq C \\ 0 \leq x_i \leq D_i \end{array} \right]$$

$$\left[ \begin{array}{l} \max \sum_{i=1}^n x_i p_i \\ \text{s.t. } \sum x_i \leq C \\ 0 \leq x_i \leq \mathbb{E}[D_i] \end{array} \right]$$