

# Optimal Stopping

• Setting - State space  $S \cup \{\text{Stopped}\}$

- Action space  $A = \{0, 1\}$

stop  $\rightarrow$   $\leftarrow$  continue

- If  $A_t = 0$ , then  $S_{t+i} = \text{'Stopped'}$  for all  $i \geq 1$

- Costs  $C_t(x, 0) = c(x)$  (continuation cost)

$C_t(x, 1) = s(x)$  (stopping cost)

$C_t(\text{Stopped}, \cdot) = 0$

• If we play this over a horizon  $T$ , then HJB eqn

$$V_t(x) = \min \left[ c(x) + \mathbb{E} \left[ V_{t+1}(X_{t+1}) \mid X_t = x, A_t = 1 \right], s(x) \right] \quad ; x \in S$$

$a=1$ , i.e., continue  $a=0$

$$V_t(\text{stopped}) = 0, \quad V_T(x) = 0 \quad \forall x$$

• 1-Step Look-Ahead Rule (1-SLA) - Stop whenever stopping now is better than stopping in next step, i.e.,  $\forall t$   $A_t = 1$  under the event

$$\mathcal{T}_t \triangleq \{ s(x_t) \leq c(x_t) + \mathbb{E}[s(x_{t+1}) \mid x_t = x, A_t = 1] \}$$

Thm (Chow & Robbins '61) - Suppose  $\mathcal{T}_t$  is monotone, i.e.  $\mathcal{T}_0 \subseteq \mathcal{T}_1 \subseteq \mathcal{T}_2 \dots$  a.s. Then 1-sla is optimal

Pf - Induction on  $t \in \{T-1, T-2, \dots, 0\}$ . For  $t=T-1$ , it's clear we should stop if  $s(x_t) \leq c(x_t) \Leftrightarrow x_t \in \mathcal{T}$ . Now if it's true for  $t \geq k$ , then at time  $t=k-1$ , if  $x_t \notin \mathcal{T}$ , then it's better to continue till  $t=k$ ; if  $x_t \in \mathcal{T}$ , then  $x_{t+1} \in \mathcal{T}$  and by induction hypothesis, it's better to stop.

Eg (Selling an asset) -  $X_1, X_2, \dots, X_T \sim F$  are iid offers for an asset with  $\text{Var}(X_i) < \infty$ ; Per day holding cost  $c$ ; also, can recall the best past offer anytime.

$$- s(x_t) = - \max_{t \leq \tau} \{X_t\}, \quad c(x_t) + E[s(x_{t+1})] = c - E[\max(X_{t+1}, M_t) | X_1, \dots, X_t]$$

$$\Rightarrow s(x_t) - c(x_t) - E[s(x_{t+1})] = \underbrace{E[(X_{t+1} - M_t)^+ | M_t]}_{\text{decreasing as fn of } M_t} - c$$

$\Rightarrow$  1-sla is optimal!

Note - Not true without recall (ie, for the 'prophet inequality')

• Corollary - If the set of 'allowed stopping points' is determined by  $X_t$ , then also we can define monotone stopping rules

Eg - Last-Success and Bruss' Odds Policy ( $q_i = 1 - p_i$ )

- Sequence of  $T$  independent events  $X_1, X_2, \dots, X_T, X_k \sim \text{Bern}(p_k)$
- Objective: Maximize probability of stopping at the 'last success' i.e.,  $R(X_T=1, A_t=0) = P[X_{t+k}=0 \forall k \geq 1]$

- Let  $\theta_i = p_i/q_i$  ('odds' for  $i^{\text{th}}$  event)  
and  $O_i = \sum_{t \geq i} \theta_t$

Thm (Bruss 2000) - The optimal stopping policy is to stop after the last time  $t$  with  $I_t=1$  and  $O_t \geq 1$

Pf - Let  $T_t = \{X_t=1, \text{ and stopping at next success is worse}\}$   
 $= \{X_t=1 \text{ and } q_{t+1}q_{t+2} \dots q_T > p_{t+1}q_{t+2} \dots q_{t+1} + \dots + q_t \dots q_{t+1} p_t\}$   
 $= \{X_t=1, 1 > \theta_{t+1} + \theta_{t+2} + \dots + \theta_T\}$

Clearly  $T_t$  is monotone  $\Rightarrow$  opt policy is stop after last  $t$  st  $\sum_{t \geq t} \theta_t \geq 1$

Corollary - For secretary problem, stop after last  $t$  st.  $H_t \geq 1$

## Optimal Stopping in Infinite Horizon Settings

- Above we assumed  $T$  was finite; can 1-sta work if  $T \rightarrow \infty$

Eg -  $X_t = X_{t+1} + Z_t$ , with  $Z_t = \pm 1$  w.p.  $1/2$  (Symmetric RW)  
 $c(x_t) = 0 \quad \forall t, \quad s(x_t) = e^{-x_t}$   
 $-s(x_t) = e^{-x_t}, \quad c'(x_t) + \mathbb{E}[s(x_{t+1})] = e^{-x_t} (e + e^{-1})/2 = 1.54e^{-x_t} > s(x_t)$   
moreover  $\tau_t = \mathbb{Z}$  which is monotone  
- However  $X_t$  is **recurrent**  $\Rightarrow$  eventually hits every  $x \in \mathbb{Z}$   
 $\Rightarrow$  opt policy is to wait forever (with cost 0)...

- However, we can avoid these pathologies with additional assumption  
Lemma - If  $S = \sup_x s(x) < \infty, C = \inf_x c(x) > 0 \Rightarrow F_T(x) \rightarrow F(x)$  as  $T \rightarrow \infty$   
Moreover a 1-sta is optimal if and only if  $\tau_t$  is monotone  
(see Weber section 6.3)

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Eg - (Sequential Prob Ratio Test) - Suppose  $X_1, X_2, \dots \sim F$ , where our priors are that  $F = F_0$  w.p.  $p_0, F = F_1$  w.p.  $p_1 = 1 - p_0$ .

- To get each additional sample, we pay cost  $\delta$
- If we stop and declare  $H_0/H_1$  true, we pay cost  $s_0/s_1$  if wrong

•  $l_0 = \frac{p_1}{p_0} = \frac{\mathbb{P}[F=F_1]}{\mathbb{P}[F=F_0]}, \quad l_n = \frac{\mathbb{P}[F=F_1 | X_1, \dots, X_n]}{\mathbb{P}[F=F_0 | X_1, \dots, X_n]} = \frac{f_1(x_n)}{f_0(x_n)} l_{n-1}$  (likelihood ratios)

• We now want to consider a 1-sta policy

- Stopping at time  $t \equiv$   
 $S(L_t) = \min \left( \frac{c_0 l_t}{1 + l_t}, c_1 \right) \Rightarrow (1 + l_t) S(L_t) = \min(c_0 l_t, c_1)$

- Continuing at  $t$  and stopping in next step

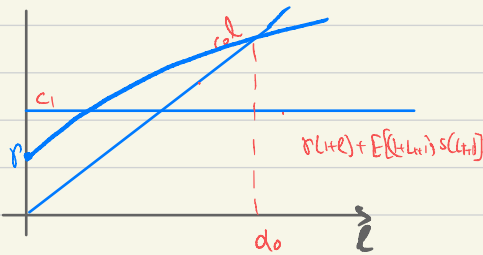
$$c(L_t) + \mathbb{E}[S(L_{t+1})|L_t] = \gamma + \frac{L_t}{1+L_t} \mathbb{E}_{x_{t+1}|F_t} [S(L_{t+1})] + \frac{1}{1+L_t} \mathbb{E}_{x_{t+1}|F_t} [S(L_{t+1})]$$

$$\begin{aligned} \Rightarrow (1+L_t)(c(L_t) + \mathbb{E}[S(L_{t+1})|L_t]) &= \gamma(1+L_t) + \int L_t f(y) S(L_t f(y)/f(y)) dy \\ &\quad + \int f(y) S(L_t f(y)/f(y)) dy \\ &= \gamma(1+L_t) + \int (1+L_t f(y)/f(y)) S(L_{t+1}) f(y) dy \\ &= \gamma(1+L_t) + \mathbb{E}_{x_{t+1}|F_t} [(1+L_{t+1}) S(L_{t+1})] \end{aligned}$$

Thus the 1-sla rule corresponds to the following:

- 1) If  $c_0 L_t \leq \gamma(1+L_t) + \mathbb{E}_{F_t} [(1+L_{t+1}) S(L_{t+1})] \Rightarrow$  Stop and declare  $\bar{F} = F_0$
- 2) If  $c_1 \leq \gamma(1+L_t) + \mathbb{E}_{F_0} [(1+L_{t+1}) S(L_{t+1})] \Rightarrow$  Stop and declare  $\bar{F} = F_1$
- 3) Else continue

- Note  $(1+l)S(l) = \min(c_0 l, c_1) \equiv$  concave, non-dec<sup>n</sup> in  $l$   
 $\Rightarrow \mathbb{E}_{F_0} [(1+L_{t+1}) S(L_{t+1})|L_t=l] = \mathbb{E}_{x_{t+1}|F_0} [\min(c_0 f(y), c_1)]$



2 intervals



3 intervals

It can be seen from the plots that the 1-sla either has a single threshold  $d_0$ , or 2 thresholds  $d_0 < d_1$ . This is clearly monotone  $\Rightarrow$  the 1-sla is optimal.